Integrity Verification of Outsourced Frequent Itemset Mining with Deterministic Guarantee

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December 10, 2013



## Data-mining-as-a-service (DMaS)

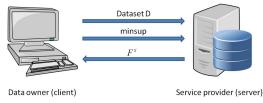
#### Data Mining as a Service:



- Weak client
- Computationally powerful service provider (e.g. cloud)
- Result integrity: are the returned mining results the same as if the computation were locally executed?

## **Outsourcing Setting**

- We focus on the problem of result integrity of outsourced *frequent itemset mining*.
- The architecture of outsourcing frequent itemset mining



Given a transaction dataset D and its correct frequent itemset mining result F, let  $F^{S}$  be the errorneous mining result that the server returns.

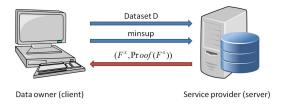
• Integrity concerns:

**Completeness** no frequent itemset is missing in  $F^S$ . **Correctness** all itemsets in  $F^S$  are frequent.

• We propose an efficient approach to catch incorrect/incomplete mining result with *100% certainty*.

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## Verification Framework



- The server constructs cryptographic proofs of the mining results.
  - We use the set intersection verification protocol[PTT11] to construct the proofs.
  - Use the proof to verify the true support of a frequent/infrequent itemset.

## Set Intersection Verification Protocol

Given a collection sets  $S = \{S_1, \ldots, S_m\}$ , an intersection result  $Y = \{y_1, \ldots, y_{\delta}\}$ ,  $Y = S_1 \cap S_2 \cap \cdots \cap S_m$  is the *correct* intersection of S if and only if:

- $(Y \subseteq S_1) \land \dots \land (Y \subseteq S_m)$  (subset condition);
- $(S_1 Y) \cap \cdots \cap (S_m Y) = \emptyset$  (completeness condition).

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[PTT11] server prepares  $\Pi(Y) = \{\mathcal{B}, \mathcal{A}, \mathcal{W}, \mathcal{C}\}$ client checks coefficients  $\mathcal{B} = \{b_{\delta}, b_{\delta-1}, \cdots, b_0\}$  of  $\mathcal{B} = \{b_0, \ldots, b_\delta\}$ polynomial  $(s + y_1)(s + y_2) \cdots (s + y_{\delta})$ are correct. accumulation values  $\mathcal{A} = \{ acc(S_i) | \forall S_i \in S \}$  $\mathcal{A}$  are correct where  $acc(S_i) = g^{\prod_{x \in S_j}(s+x)}$  $e(\prod_{k=0}^{|Y|} (g^{s^k})^{b_k}, W_i)$ subset witness  $\mathcal{W} = \{W_i | \forall S_i \in \mathcal{S}\}$  $\stackrel{?}{=} e(acc(S_i), g)$ where  $W_i = g^{P_j(s)}$ ,  $P_i(s) = \prod_{x \in S_i - Y} (x + s)$ for  $j = 1, \cdots, m$  $\overline{\prod_{i=1}^m e(W_j, C_j)}$ completeness witness  $C = \{C_i | \forall S_i \in S\}$  $\stackrel{?}{=} e(g,g)$ for each set  $S_i \in S$ ,  $C_i = g^{q_j(s)}$ s.t.  $q_1(s)P_1(s) + q_2(s)P_2(s) + \cdots + q_m(s)P_m(s) = 1$ 

## **Basic Solution**

Given a dataset D that contains n unique items, the client does the following:

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- Build the *item-based inverted index E<sup>1</sup>* that consists of n inverted lists {L<sub>1</sub>,..., L<sub>n</sub>}.
- **2** Construct the Merkle hash tree  $\mathcal{T}$  of the inverted index.
  - Leaf  $I_j$  is assigned  $h_j = hash(acc(L_j)^{(s+j)})$ .
  - Internal node v with children  $c_1, \ldots, c_k$  is assigned  $h_v = hash(h_{c_1}|| \ldots ||h_{c_k}).$

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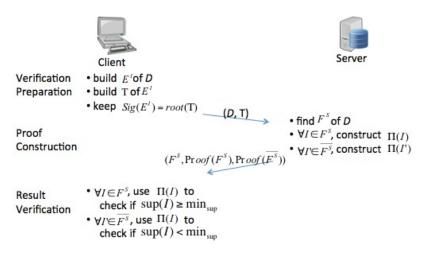
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Mapping to the set intersection verification problem Verifying whether any itemset I is included in a set of transactions T' is equivalent to verifying whether T' is the correct intersection of the inverted lists of all items in I.

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## **Basic Solution**



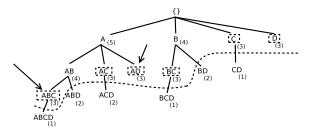
Drawbacks

- Total number of proofs is  $2^n 1$ .
- Too much overhead.

## Verification Optimization

Maximal frequent itemset (MFI) A subset of  $F^{S}$  s.t. for each itemset  $I \in MFI$ , there does not exist any itemset  $I' \in F^{S}$  s.t.  $I \subseteq I'$ .

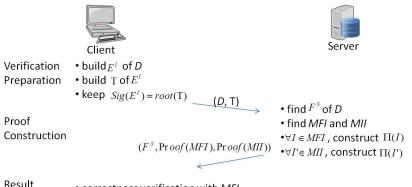
Minimal infrequent itemset (MII) A set of itemsets that do not appear in  $F^{S}$  s.t. for each itemset  $I \in MII$ , there does not exist any itemset  $I' \notin F^{S}$  s.t.  $I' \subseteq I$ .



(Itemsets in dotted rectangles are maximal frequent itemsets.)

Advantage  $|MFI| + |MII| \ll |F^{S}| + |\overline{F^{S}}|$ 

## **Optimized Solution**



Verification

• correctness verification with *MFI* 

• completeness verification with MII

Security Analysis Our optimized solution provides the same security guarantee as the basic solution.

Proof construction at server side  $O(Mlog^3M + n^{\epsilon}logn)$ •  $M = \sum_{I \in MFI \cup MII} \sum_{i \in I} |L_i|$ • n is the number of unique items of D. •  $\epsilon \in (0, 1)$ Verification at client side O(N + F)•  $N = \sum_{I \in MFI \cup MII} |I|$ •  $F = \sum_{I \in MFI \cup MII} sup(I)$ 

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## Experiments

#### Environment

Language C++ Testbed Macbook Pro, 2.4*GHz* CPU, 4 GB memory

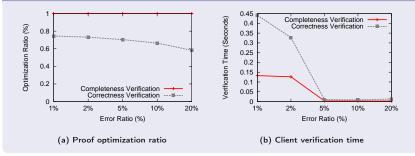
•	Dataset	# of	# of	Avg. trans.	min <sub>sup</sub>	# of freq.
		trans.	items	length		itemsets
	$S_1$	10 <sup>3</sup>	49	10	250	36
	<i>S</i> <sub>2</sub>	10 <sup>4</sup>	49	10	250	3854
	$S_3$	10 <sup>5</sup>	49	10	250	149744
	<i>S</i> <sub>4</sub>	10 <sup>6</sup>	49	10	250	3074610
	R	500	100	2.4	5	97

• Simulation of malicious actions

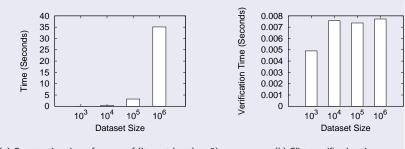
Error ratio r = 1%, 2%, 5%, 10%, 20%Incomplete Randomly delete *r* percent mining result. Incorrect Randomly insert *r* percent infrequent itemsets.

# Proof Optimization Ratio & Verification Time

#### **Optimization Ratio & Verification Time (***R* **dataset)**



#### Scalability (error ratio=1%)



(a) Construction time of one proof (itemset length = 3)

(b) Client verification time

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## Thank you!

Questions?

## **Related Work**

#### Verifiable Computation

• [Bab85, GMR89, PRV12, GGP10] the expensive pre-processing phase is amortized over the future executions.

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- [Bab85, GMR89, PRV12, GGP10] the expensive pre-processing phase is amortized over the future executions.
- Integrity Verification of Database-as-a-Service (DaS)
  - [PJRT05, Sio05, XWYM07] provide assurance for SQL query results.

## **Related Work**

#### Verifiable Computation

- [Bab85, GMR89, PRV12, GGP10] the expensive pre-processing phase is amortized over the future executions.
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#### Integrity Verification of DMaS

- [WCH<sup>+</sup>09, DLW13] only provide probabilistic result integrity guarantee.
- [LWM<sup>+</sup>12, RHPH13] focus on other mining tasks (outlier detection, clustering)

#### Comparison on $S_1$ dataset

min <sub>sup</sub>	# of Freq.	Client side	Server side	
	Itemsets	Verify	Proof prep.	mining
402	10	0.000164	24.72	0.03707
203	50	0.001358	266.985	0.08984
157	99	0.00332	572.591	0.1355

(time measured in seconds)

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